
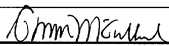


09/21/00
09/666456
U.S. PTO

UTILITY PATENT APPLICATION TRANSMITTAL <small>(Only for new nonprovisional applications under 37 CFR 1.53(b))</small>	Attorney Docket No.	197529US2
	First Inventor or Application Identifier	Takashi MIYAMORI
	Title	PROCESSOR CORE FOR USING EXTERNAL EXTENDED ARITHMETIC UNIT EFFICIENTLY AND PROCESSOR INCORPORATING THE SAME

APPLICATION ELEMENTS <small>See MPEP chapter 600 concerning utility patent application contents</small>	ADDRESS TO: Assistant Commissioner for Patents Box Patent Application Washington, DC 20231
1. <input checked="" type="checkbox"/> Fee Transmittal Form (e.g. PTO/SB/17) <small>(Submit an original and a duplicate for fee processing)</small>	ACCOMPANYING APPLICATION PARTS 6. <input type="checkbox"/> Assignment Papers (cover sheet & document(s)) 7. <input type="checkbox"/> 37 C.F.R. §3.73(b) Statement <input type="checkbox"/> Power of Attorney <small>(when there is an assignee)</small> 8. <input type="checkbox"/> English Translation Document (if applicable) 9. <input checked="" type="checkbox"/> Information Disclosure Statement (IDS)/PTO-1449 <input checked="" type="checkbox"/> Copies of IDS Citations (1) 10. <input type="checkbox"/> Preliminary Amendment 11. <input checked="" type="checkbox"/> White Advance Serial No. Postcard 12. <input type="checkbox"/> Small Entity Statement(s) <input type="checkbox"/> Statement filed in prior application. Status still proper and desired. 13. <input checked="" type="checkbox"/> Certified Copy of Priority Document(s) (1) <small>(if foreign priority is claimed)</small> 14. <input checked="" type="checkbox"/> Other: Notice of Priority
2. <input checked="" type="checkbox"/> Specification Total Pages 29	
3. <input checked="" type="checkbox"/> Drawing(s) (35 U.S.C. 113) Total Sheets 8 <small>(Formals)</small>	
4. <input checked="" type="checkbox"/> Oath or Declaration Total Pages 2 a. <input checked="" type="checkbox"/> Newly executed (original) b. <input type="checkbox"/> Copy from a prior application (37 C.F.R. §1.63(d)) <small>(for continuation/divisional with box 15 completed)</small> i. <input type="checkbox"/> DELETION OF INVENTOR(S) <small>Signed statement attached deleting inventor(s) named in the prior application, see 37 C.F.R. §1.63(d)(2) and 1.33(b)</small>	
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Docket No. 197529US2

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

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SERIAL NO: New Application

FILING DATE: Herewith

FOR: PROCESSOR CORE FOR USING EXTERNAL EXTENDED ARITHMETIC UNIT EFFICIENTLY AND
PROCESSOR INCORPORATING THE SAME

FEE TRANSMITTAL

ASSISTANT COMMISSIONER FOR PATENTS
WASHINGTON, D.C. 20231

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TOTAL CLAIMS	19 - 20 =	0	× \$18 =	\$0.00
INDEPENDENT CLAIMS	3 - 3 =	0	× \$78 =	\$0.00
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PROCESSOR CORE FOR USING EXTERNAL EXTENDED ARITHMETIC UNIT
EFFICIENTLY AND PROCESSOR INCORPORATING THE SAME

BACKGROUND OF THE INVENTION

5

1. Field of the Invention

The present invention relates to a processor core for using an external extended arithmetic unit efficiently and a processor incorporating the same, and more particularly to a technique to improve processing efficiency in a processor core that makes an external extended arithmetic unit execute a part of an arithmetic operation instruction by using the external extended arithmetic unit efficiently.

15 2. Description of the Background Art

Fig. 1 is a block diagram showing an example of a basic arrangement of a conventional processor.

The processor includes a processor core 70 and a data memory 80. The processor core 70 comprises a fetch PC (FPC) 20 701 for fetching an instruction, an address incrementer 702 for incrementing an address in an instruction memory, an instruction memory 703 for storing an instruction, an instruction register 704 for retaining an instruction read out from the instruction memory 703, a register file 705, a 25 decode circuit 706 for decoding an instruction, a source (1) register 707 and a source (2) register 708, an execute stage control register 710, an arithmetic circuit 711 for executing an arithmetic operation, a pipeline control circuit 712 for controlling pipeline processing in the processor, a memory 30 stage data register 713, a memory stage control register 715, a register write stage data register 716, a register write stage control register 717, and an instruction validating register 718.

The processor core 70 executes an instruction by, for 35 example, a 5-stage pipeline structure. In other words, in the processor core 70, the pipeline stages include a pipeline

stage 1 (instruction fetch stage), a pipeline stage 2 (register read stage), a pipeline stage 3 (execute stage), a pipeline stage 4 (memory stage), and a pipeline stage 5 (register write stage).

5 The instruction fetch stage is a stage where an instruction is read out from the instruction memory 703, and the register read stage is a stage where a register specified by the instruction read out in the instruction fetch stage is read out from the register file 705 while the instruction
10 is decoded by the decode circuit 706. The execute stage is a stage where the instruction is executed by the arithmetic circuit 711 in accordance with a value in the register read out in the register read stage and the decode information of the instruction, and the memory stage is a stage where an access
15 is made to the data memory 80 in case that the instruction is a memory load or store instruction. The register write stage is a stage where an execution result in the execute stage or load data in case of a memory load instruction is written into the register file 705.

20 However, the conventional processor core merely executes a pre-defined instruction, and is not provided with an extended function for connecting an arithmetic unit to its exterior. In other words, because the conventional processor is not provided with mechanism or an interface signal for connecting
25 an arithmetic unit to its exterior efficiently, the system performance cannot be improved by connecting an arithmetic unit suitable for application systems to the processor core efficiently.

30 Pre-integrated an arithmetic unit suitable to application systems, for example, a production-sum arithmetic circuit, in the processor core may eliminate the above problem. However, not all the application systems use the production-sum arithmetic circuit. Hence, incorporating the production-sum arithmetic circuit in every processor core may produce useless
35 hardware, thereby causing the cost to be increased unnecessarily.

Fig. 2 is a block diagram showing a second arrangement of the conventional processor.

The processor of the arrangement shown in Fig. 2 is connected to a coprocessor at its exterior. The coprocessor receives an instruction directed to the coprocessor from the processor core, and executes the same.

The coprocessor includes in its interior a coprocessor register file 705b, a coprocessor source (1) register 707b and a coprocessor source (2) register 708b, a coprocessor arithmetic circuit 711b, and a pipeline register 720b.

The coprocessor executes an arithmetic operation specified by the instruction directed to the coprocessor by reading out a value in each register in the coprocessor register file 705b specified by the above instruction, and using the read out values as input data to the coprocessor arithmetic circuit 711b.

With the processor core of the arrangement shown in Fig. 2, by connecting the coprocessor to its exterior, the function can be extended, but data used in the arithmetic operation carried out by the coprocessor arithmetic circuit 711b is still limited to the content of the coprocessor register file 705b in the coprocessor. For this reason, the coprocessor arithmetic circuit 711b is not allowed to execute an arithmetic operation by directly using the content of the processor register file 705 in the processor core.

As has been discussed, the conventional processor core is not provided with an interface function for connecting an arithmetic unit to its exterior efficiently, and there has been a need for an extended function for executing an arithmetic operation efficiently by using an external arithmetic unit connected to the processor core.

SUMMARY OF THE INVENTION

It is therefore an object of the present invention to provide a processor core for connecting an arithmetic unit

to its external efficiently, so that the system performance can be improved drastically without increasing the size thereof by connecting the arithmetic unit to its exterior to furnish an efficient interface function therebetween.

5 To achieve the object, an aspect of the invention provides a processor, comprising: a processor core; a data memory accessed by the processor core; and an extended arithmetic unit, connected to an exterior of the processor core, for processing a particular instruction, the extended arithmetic
10 unit executing an arithmetic operation by using arithmetic operation data retained in a register file in the processor core and outputting a result of an arithmetic operation directly to the processor core, the processor core saving the result of the arithmetic operation executed by the extended
15 arithmetic unit and inputted therefrom in the register file in the processor core.

Another aspect of the invention provides a processor, comprising: a processor core; a data memory accessed by the processor core; and an extended arithmetic unit, connected
20 to an exterior of the processor core, for processing a particular instruction, the processor core, at least including: an instruction memory for storing an instruction to be executed; an instruction decode unit for reading out an instruction from the instruction memory to decode the
25 instruction, in case that the instruction decoded is an extended arithmetic unit control instruction that should be executed by the extended arithmetic unit connected to the exterior of the processor core, the instruction decode unit also outputting at least an instruction code of the extended
30 arithmetic unit control instruction to the extended arithmetic unit; a register file for retaining arithmetic operation data of an arithmetic operation that should be executed by the instruction decoded, in case that the arithmetic operation data is data of the extended arithmetic unit control
35 instruction, the register file also outputting the arithmetic operation data to the extended arithmetic unit; a first

operational section for executing the instruction decoded; and an extended arithmetic unit, at least including, a second operational section for executing an arithmetic operation specified by the extended arithmetic unit control instruction by using the arithmetic operation data retained in the register, and outputting an execution result of the arithmetic operation to the processor core.

10 Preferably, in case that the instruction decoded is the extended arithmetic unit control instruction, the processor core outputs to the extended arithmetic unit at least an instruction code that specifies an action involved in an arithmetic operation in the extended arithmetic unit and an instruction valid signal that indicates the instruction code is valid.

20 Preferably, the processor core includes a pipeline control unit for controlling pipeline processing in an interior of the processor core and in the extended arithmetic unit.

Preferably, the pipeline control unit outputs to the extended arithmetic unit a pipeline flush signal for abandoning execution of an instruction outputted thereto.

Preferably, the extended arithmetic unit outputs to the
35 processor core an arithmetic operation result invalidating
signal that invalidates an execution result of an arithmetic

operation executed therein.

Preferably, the data memory receives from the extended arithmetic unit at least one of an address in memory access, data, a write control signal for controlling data writing, and a read control signal for controlling data reading; reads out the data from a region specified by the address and outputs the data to the extended arithmetic unit in case that data reading is carried out because the read control signal is asserted; and writes the data inputted from the extended arithmetic unit into a region specified by the address in case that data writing is carried out because the write control signal is asserted.

Preferably, the extended arithmetic unit includes: a plurality of arithmetic circuits; a first pipeline register for storing a processing result by an arithmetic circuit in a preceding stage at a rising of a following clock; and a second pipeline register for storing a processing result by an arithmetic circuit in a succeeding stage at the rising of the following clock.

Still another aspect of the present invention provides a processor core connected to an extended arithmetic unit for processing a particular instruction to an exterior thereof, comprising: an instruction memory for storing an instruction to be executed; an instruction decode unit for reading out an instruction from the instruction memory to decode the instruction, in case that the instruction decoded is an extended arithmetic unit control instruction that should be executed by the extended arithmetic unit connected to the exterior of the processor core, the instruction decode unit also outputting at least an instruction code of the extended arithmetic unit control instruction to the extended arithmetic unit; a register file for retaining arithmetic operation data of an arithmetic operation that should be executed by the instruction decoded, and in case that the arithmetic operation data is data for the extended arithmetic unit control instruction, the register file also outputting the arithmetic

operation data to the extended arithmetic unit and storing a result of an arithmetic operation executed in the extended arithmetic unit.

Preferably, in case that the instruction decoded is the
5 extended arithmetic unit control instruction, the instruction decode unit outputs to the extended arithmetic unit at least an instruction code that specifies an action involved in an arithmetic operation by the extended arithmetic unit and an instruction valid signal that indicates the instruction code
10 is valid.

Other features and advantage of the present invention will become apparent from the following description taken in conjunction with the accompanying drawings.

15 BRIEF DESCRIPTION OF DRAWINGS

The accompanying drawings, which are incorporated in and constitute a part of the specification, illustrate presently preferred embodiments of the invention, and together
20 with the general description given above and the detailed description of the preferred embodiments given below, serve to explain the principles of the invention.

Fig. 1 is a block diagram showing an example of a basic arrangement of a conventional processor core;

25 Fig. 2 is a block diagram showing an example of a basic arrangement of the conventional processor core to which a coprocessor is connected;

Fig. 3 is a block diagram showing a basic arrangement of a first embodiment of a processor core and a processor
30 incorporating the same according to the present invention;

Fig. 4 is a view showing an example of bit patterns of instructions executed by the processor of the first embodiment;

Fig. 5 is a view showing an input and an output of a decode circuit 106 of Fig. 3;

35 Fig. 6 is a view showing an input and an output of an arithmetic circuit 111 of Fig. 3;

Fig. 7 is a block diagram showing a basic arrangement of a second embodiment of a processor core and a processor incorporating the same according to the present invention; and

Fig. 8 is a block diagram showing a basic arrangement of a third embodiment of a processor core and a processor incorporating the same according to the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

The following description will describe in detail embodiments of a processor core and a processor incorporating the same in the present invention with reference to Figs. 3 to 8.

First Embodiment

The following description will describe in detail a first embodiment of a processor core and a processor incorporating the same according to the present invention with reference to Figs. 3 to 6.

The first embodiment provides a function for improving the system's processing efficiency by providing a processor that executes pipeline processing with an interface for connecting an arithmetic unit to the exterior of the processor core so that the arithmetic unit executes an arithmetic operation.

Fig. 3 is a block diagram showing an arrangement of a processor in accordance with the first embodiment of the present invention.

The processor of the first embodiment includes a processor core 10, an extended arithmetic unit 20 connected to the exterior of the processor core 10, and a data memory 30.

The processor core 10 comprises a fetch program counter (FPC) 101, an address incrementer 102, an instruction memory 103, an instruction register 104, a register file 105, a decode

circuit 106, a source (1) register 107 and a source (2) register 108, an execute stage control register 110, an arithmetic circuit 111, a pipeline control circuit 112, a memory stage data register 113, a memory stage control register 115, a
5 register write stage data register 116, a register write stage control register 117, and an instruction validating register 118.

The processor core 10 executes an instruction by, for example, a 5-stage pipeline structure. In other words, the
10 processor core 10 includes the following pipeline stages: a pipeline stage 1 (instruction fetch stage) as a stage where an instruction is read out from the instruction memory 103; a pipeline stage 2 (register read stage) as a stage where a register specified by the instruction read out in the
15 instruction fetch stage is read out from the register file 105 while the instruction is decoded by the decode circuit 106; a pipeline stage 3 (execute stage) as a stage where the instruction is executed by the arithmetic circuit 111 by using a value in the register read out in the register read stage
20 and the decode information of the instruction; a pipeline stage 4 (memory stage) as a stage where an access is made to the data memory 30 in case that the instruction is a memory load or store instruction; and a pipeline stage 5 (register write stage) as a stage where an execution result in the execute
25 stage or load data in case of a memory load instruction is written into the register file 105.

Typical instructions for the processor core 10 of the first embodiment include an add instruction (add), a data memory load instruction (load), a data memory store instruction
30 (store), a jump instruction (jump), and an extended instruction (ext), and each instruction is 32-bit long, for example.

Fig. 4 is a view showing an example of a bit pattern of each instruction used in the processor core, and actions involved with each instruction will be explained with reference
35 to Fig. 4.

The add instruction (add) is a 32-bit instruction whose

19th to 16th bits show "0000", and a value in a register specified by the src (source) 1 field is added to a value in a register specified by the src2 field, and the addition result is written into a register specified by the dst (destination) field.

5 The data memory load instruction (load) is an instruction whose 19th to 16th bits show "0001", and data is read out from the data memory 30 by using a value in a register specified by the src1 field as an address in the data memory 30, and the readout data is written into a register specified by the
10 dst field.

 The data memory store instruction (store) is an instruction whose 19th to 16th bits show "0010", and a value in a register specified by the src2 field is written into the data memory 30 by using a value in a register specified by
15 the src1 field as an address in the data memory 30.

 The jump instruction (jump) is an instruction whose 19th to 16th bits show "0011", and a value in a register specified by the src1 field is read out, and a jump is taken by using the readout value as a jump destination address.

20 The extended instruction (ext) is an instruction whose 19th to 16th bits show "0100", and values in registers respectively specified by the src1 field and src2 field are read out to be outputted to an external extended arithmetic unit 20. The extended arithmetic unit 20 outputs an arithmetic
25 operation result to the processor core 10, and the result is written into a register specified by the dst field.

 Next, the following description will describe the basic operations of the processor core of the first embodiment in each of the foregoing pipeline stages.

30 (1) Operations in the Pipeline Stage 1

 In the pipeline stage 1 (instruction fetch stage), the following operations are carried out repetitively. That is, an instruction specified by a value in the fetch PC 101 is read out from the instruction memory 103. The instruction
35 read out from the instruction memory 103 is written into the instruction register 104 and a value in the instruction

validating register 118 is set to 1 at the rising of the following clock. The number of bytes of an instruction word (four in case of a 4-byte instruction word) is added to the value in the fetch PC 101 in the address incrementer 102, and the result is stored in the fetch PC 101 at the rising of the following clock, after which the following instruction is read out from the instruction memory 103.

(2) Operations in the Pipeline Stage 2

In the pipeline stage 2 (register read stage), a value in a register specified by the instruction in the instruction register 104 is read out from the register file 105. In case of the instruction shown in Fig. 4, the 27th to 24th bits of the instruction are assigned as the src1 field, and the 23rd to 20th bits are assigned as the src2 field, and a value in a register specified by each set of bits is read out from the register file 105. The values in the registers read out from the register file 105 are stored respectively into the source (1) register 107 and source (2) register 108 at the rising edge of the following clock. As it will be described below, the values in the registers read out from the register file 105 are also stored in an execute stage register 201 in the extended arithmetic unit, and used as input data for an arithmetic operation executed by the extended arithmetic unit.

The decode circuit 106 decodes the instruction in the instruction register 104 when a value "1" is set in the instruction validating register 118 and the instruction stored in the instruction register 104 is valid. As a result of this decoding, the decode circuit 106 outputs control information of the arithmetic circuit 111, the number and valid information of the destination register, and pipeline flush information in case of a jump instruction to the execute stage control register 110. Each kind of the foregoing information is stored into the execute stage control register 110 at the rising of the following clock.

Further, the operation of the decode circuit 106 will be explained more in detail.

data from the arithmetic circuit 111 includes the contents of the source (1) register 107 and source (2) register 108. Also, when the control information of the arithmetic circuit 111 indicates a value "10" (ext), a selector 109 does not select an output from the arithmetic circuit 111 but an input from the extended arithmetic unit 20, which is delivered to the memory stage.

The arithmetic operation result (output data) from the arithmetic circuit 111 is stored in the memory stage data register 113, and in case that the instruction is a data load instruction or a data store instruction demanding an access to the data memory 30, an output from the arithmetic circuit 111 is outputted as an address in the data memory 30. In case of a data store instruction, a value from the selector 109 is outputted as data to be stored in the data memory 30.

The number and valid information of the destination register in the execute stage control register 110 are stored in the memory stage control register 115 at the rising of the following clock.

When the pipeline flush information is validated by a jump instruction, the pipe line control circuit 112 asserts a pipeline flush signal, and invalidates the valid information of the instruction validating register 118 and execute stage control register 110 at the rising of the following clock, so that neither the decode stage nor execute stage will be operated in the following clock.

In addition, when the pipeline flush signal is activated, the pipeline control circuit 112 effects a control in such a manner that the following value in the fetch PC 101 is the jump destination address of an output from the arithmetic circuit 111.

At the same time, the pipeline control circuit 112 outputs the pipeline flush signal to the extended arithmetic unit 20 so as to disable the execute stage register 201 in the extended arithmetic unit 20.

When a pipeline stall (stop) request is issued from the

data memory 30 or a pipeline stop request is issued from the extended arithmetic unit 20, the pipeline control circuit 112 outputs a control signal to each module so that a value in each of the fetch PC 101, instruction register 104, instruction
5 validating register 118, source (1) register 107, source (2) register 108, and execute stage control register 110 will not be updated at the rising of the following clock.

In case of a pipeline stall request from the data memory 30, the pipeline control circuit 112 outputs the control signal
10 also to the memory stage control register 115, so that a value in the memory stage control register 115 will not be updated either.

At the same time, the pipeline control circuit 112 outputs a pipeline stop signal to the extended arithmetic unit 20 to
15 effect a control in such a manner that a value in the execute stage register 201 in the extended arithmetic unit 20 will not be updated.

(4) Operations in the Pipeline Stage 4

In the pipeline stage 4 (memory stage), in case that
20 the instruction is a data load instruction, data is read out from the data memory 30 and stored in the register write stage data register 116 in the following clock. In case that the instruction is any other kind of instruction, a value in the memory stage data register 113 is stored in the register write
25 stage data register 116 in the following clock. In addition, information of the memory stage control register 115 is stored in the register write stage control register 117 in the following clock.

(5) Operations in the Pipeline Stage 5

In the pipeline stage 5 (register write stage), when
30 the valid information in the register write stage control register 117 shows valid, a control is effected in such a manner that a value in the register write stage data register 116 is written into a register in the register file 105 specified
35 by the destination register number in the register write stage control register 117.

Next, the following description will describe an arrangement and operations of the extended arithmetic unit 20 of the first embodiment.

5 The extended arithmetic unit 20 is connected to the processor core 10, and receives from the processor core 10 a value in a register corresponding to each of the src1 and src2 fields in the instruction read out from the register file 105 in the decode stage thereof.

Also, the extended arithmetic unit 20 receives extended
10 instruction information (code 16) corresponding to the lower 16 bits in the instruction from the decode circuit 106, and extend instruction valid information that indicates a value "1" when the instruction in the decode stage is an extended instruction ext. Each kind of the foregoing information is
15 stored in the execute stage register 201 at the rising of the following clock unless the pipeline flush signal or the pipeline stall signal from the processor core 10 is asserted. On the other hand, when the pipeline stall signal from the processor core 10 is asserted, a value in the execute stage
20 register 201 is not updated.

When the pipeline flush signal from the processor core 10 is asserted, the extended instruction valid information is forcibly invalidated (0) at the rising of the following clock, so that the extended arithmetic unit 20 will not carry
25 out an arithmetic operation.

A stage (1) arithmetic circuit 202 is a unit that carries out an arithmetic operation in a stage corresponding to the memory stage in the processor core 10. The stage (1) arithmetic circuit 202 is operated by decoding the extended instruction
30 information code 16 inputted from the processor core 10 by means of a control circuit 204. The stage (1) arithmetic circuit 202 can use a part of the extended instruction information as input data of an arithmetic operation.

An arithmetic operation result by the stage (1)
35 arithmetic circuit 202 is stored in the pipeline (1) register 205 at the rising of the following clock. A result in a pipeline

(1) register 205 can be used as input data to a following stage
(2) arithmetic circuit 203. A result in the stage (2) arithmetic circuit 203 is stored in a pipeline (2) register 206 at the rising of the following clock.

5 The following is an example of an extended instruction sequence in the first embodiment:

1. ext dst,src1,src2,0x0001 acc <- src1*src2;dst <- src1
2. ext dst,src1,src2,0x0002 acc <- acc+src1*src2;dst <- src1
3. ext dst,src1,src2,0x0003 dst <- acc

10 In order to execute these instructions 1 to 3, the stage (1) arithmetic circuit 202 in the extended arithmetic unit 20 carries out a multiplication (instruction 1), and the stage (2) arithmetic circuit 203 adds the multiplication result of the stage (1) arithmetic circuit 202 inputted from the pipeline
15 (1) register 205 to a value in the pipeline (2) register 206 (instruction 2). A result in the stage (2) arithmetic circuit 203 is stored again in the pipeline (2) register 206 (in this case, acc register).

With the instruction 1, the stage (1) arithmetic circuit
20 202 carries out a multiplication of the inputs from the processor core 10: source (1) data (src1) and source (2) data (src2), and a result is stored in the pipeline (1) register 205 at the rising of the following clock. The stage (2) arithmetic circuit 203 remains inactive, and the
25 multiplication result stored in the pipeline (1) register 205 is stored in the pipeline (2) register 206 at the rising of the following clock. In addition, the stage (1) arithmetic circuit 202 outputs the source (1) data to the processor core 10 as a result of the instruction execution. This corresponds
30 to the execute stage in the processor core 10, and the processor core 10 writes the result of the instruction 1 in the register file 105 at the processor core 10 side in the register write stage.

With the instruction 2, the operations up to the stage
35 1 are the same as those with the instruction 1. The stage (2) arithmetic circuit 203 adds a result in the pipeline (1)

register 205 to a result in the pipeline (2) register 206,
and stores the addition result in the pipeline (2) register
206. The stage (1) arithmetic circuit 202 outputs the source
(1) data to the processor core 10 as a result of the instruction
5 execution. This corresponds to the execute stage in the
processor core 10, and the processor core 10 writes the result
of the instruction 2 in the register file 105 at the processor
core 10 side in the register write stage.

With the instruction 3, the pipeline (2) register 206
10 is read out, and the stage (1) arithmetic circuit 202 outputs
the source (1) data to the processor core 10 as a result of
the instruction. This corresponds to the execute stage in
the processor core 10, and the processor core 10 writes the
result of the instruction 3 in the register file 105 in the
15 register write stage.

The result of either of the instructions 1 or 2 is stored
in the pipeline (2) register 206. Thus, when the instruction
3 is stored in the execute stage register 201, the control
circuit 204 checks whether the instruction 1 or 2 is stored
20 in the pipeline (1) register 205. In this case, the pipeline
stop signal (pipeline stall signal) to the processor core 10
is asserted for one cycle to stop the pipeline processing at
the processor core 10 side and read out the result of the
instruction 1 or 2 in the following clock cycle, thereby
25 outputting the result to the processor core 10 through the
stage (1) arithmetic circuit 202.

In other words, the control circuit 204 checks an
instruction that demands more than one cycle when executed
at the extended arithmetic unit 20 side, and asserts the
30 pipeline stop signal to the processor core 10 side, so that
the pipeline processing at the processor core 10 side is
suspended until the timing at which the arithmetic operation
result can be outputted to the processor core 10 side.

As has been discussed, according to the first embodiment,
35 the processor core 10 is provided with an extended instruction
for controlling the extended arithmetic unit 20, and the

control circuit for executing the extended instruction and the interface circuit for controlling the extended arithmetic unit 20 and transferring data are incorporated therein. Consequently, the extended arithmetic unit 20 can be connected to the exterior of the processor core 10 efficiently, and the extended arithmetic unit 20 can execute an arithmetic operation by using the content of the register file 105 at the processor core 10 side as an input data, and output an arithmetic operation result directly to the processor core 10.

For example, assume that a product-sum operational circuit having a latency of two cycles and a throughput of one cycle is provided as the extended arithmetic unit 20, and it takes for the processor 10 five cycles and one cycle to execute a multiplication instruction and an add instruction, respectively. Then, ten times of product-sum operations can be executed in 11 cycles by using the product-sum operational circuit serving as the extended arithmetic unit 20. In contrast, in case of a combination of the multiplication instruction and add instruction by the processor core 10, $(5+1) \times 10 = 60$ cycles are necessary.

The product-sum operational circuit may be pre-incorporated in the processor core 10, but a product-sum operation is not used at all in many application systems, and the product-sum arithmetic unit is useless hardware for these application systems, thereby causing the cost to be increased unnecessarily.

On the other hand, according to the first embodiment, besides the above product-sum operation example case, by connecting various kinds of arithmetic units suitable for application systems to the exterior of the processor core 10 and downsizing the arrangement of the processor core 10 to the minimum, the system performance can be improved by executing processing efficiently between the processor core 10 and extended arithmetic unit 20 while saving the manufacturing costs of the processor core 10.

In the above explanation, the extended arithmetic unit

20 is of a 2-stage pipeline structure composed of the stage
(1) arithmetic circuit 202 and stage (2) arithmetic circuit
203. However, the arrangement of the extended arithmetic unit
20 is not limited to the foregoing, and the extended arithmetic
5 unit 20 can have three or more stages as the pipeline stages.

In addition, the case of reading out the result of an
instruction that demands more than one stage when executed
was explained as an example of a pipeline stall request
(asserting a pipeline stop signal) to the processor core 10.
10 However, the example is not limited to the foregoing. For
example, when pipeline processing is not carried out in the
interior and an instruction demands more than one cycle when
executed, the execution of the following instruction may be
suspended until the currently executed instruction is
15 completed. This can be realized by managing the state of
instruction execution by the control circuit 204.

Unlike the coprocessor connected to the conventional
processor shown in Fig. 2, the extended arithmetic unit 20
of the first embodiment carries out an arithmetic operation
20 by using data in the register file 105 in the processor core
10, and the processor core 10 receives an arithmetic operation
result outputted from the extended arithmetic unit 20 and saves
the same in the register file 105 therein. Consequently, the
system's processing efficiency can be improved more readily.

25

Second Embodiment

Next, with reference to Fig. 7, the following description
will describe in detail a second embodiment of a processor
core and a processor incorporating the same according to the
30 present invention as to differences from the first embodiment
alone.

Compared with the first embodiment shown in Fig. 3, the
second embodiment is an embodiment additionally provided with
a function of directly accessing the data memory.

35 Fig. 7 is a block diagram showing an example of an
arrangement of a processor in accordance with the second

embodiment of the present invention. Like members with respect to Fig. 3 are designated by like reference characters and the explanation of these members is not repeated for ease of explanation.

5 In the processor of the second embodiment, the processor core 10 is of the same arrangement as that shown in Fig. 3, but different in that a stage (1) arithmetic circuit 402 in an extended arithmetic unit 40 can directly access a data memory 31 by outputting an address signal, a read/write control signal, and a write data signal to the data memory 31.

The data memory 31 selects the address signal, read/write control signal, write data signal from the processor core 10 by means of selectors 31a and 31b, and outputs the selected signal to a RAM module, whereby access requests from the processor core 10 and extended arithmetic unit 40 are arbitrated.

Data read out from the data memory 31 is outputted to the extended arithmetic unit 40, and either this data or an output of a stage (2) arithmetic circuit 403 is selected and stored in a pipeline (2) register 406 at the rising of the following clock.

In the extended arithmetic unit 40 of the second embodiment, a stage (3) arithmetic circuit 407 uses an output from the pipeline (2) register 406 as input data of an arithmetic operation. In this manner, a result read out from the data memory 31 can be used as arithmetic operation input data in the extended arithmetic unit 40.

According to the second embodiment, because data directly read out from the data memory 31 can be used as arithmetic operation input data in the extended arithmetic unit 40, an instruction demanding memory access can be executed by the extended arithmetic unit 40.

Third Embodiment

35 Next, with reference to Fig. 8, the following description will describe in detail a third embodiment of a processor core

and a processor incorporating the same according to the present invention as to differences from the foregoing embodiments alone.

Fig. 8 is a block diagram showing an example of an arrangement of a processor in accordance with the third embodiment of the present invention. Like members with respect to Fig. 3 are designated by like reference characters and the explanation of these members is not repeated for ease of explanation.

Compared with the first embodiment shown in Fig. 3, the processor of the third embodiment is an embodiment such that, by outputting destination valid information from the extended arithmetic unit 20 to the processor core 10, an arithmetic operation result in the execute stage in the processor core 10 is invalidated by the processor core 10 so as not to be written into a register in the register write stage. The destination valid information is the information that shows whether an arithmetic operation result should be written into the destination register in the processor core 10, and for example, can be realized by a bit that enables/disables the destination register in the processor core 10.

In the third embodiment, the extend instruction shown in the first embodiment is changed as follows:

```
1. ext src1,src2,0x0001    acc <- src1*src2;
2. ext src1,src2,0x0002    acc <- acc+src1*src2;
3. ext dst,src1,src2,0x0003 dst <- acc
```

The register acc is a register in the extended arithmetic unit 20, and results of the instruction 1 and 2 are not written back into the register file 105 in the processor core 10.

For this reason, a control circuit 204a in the extended arithmetic unit 20 invalidates the destination valid information and outputs the same to the processor core 10. The destination register valid information of the execute stage control register 110 in the processor core 10 indicates a value of invalid (0) in this case, which is stored in a memory stage control register 115a at the rising of the following clock.

With the instruction 3, the result of an arithmetic operation executed by the extended arithmetic unit 20 is written back into the register file 105 in the processor core 10. For this reason, the control circuit 204a in the extended arithmetic unit 20 validates the destination valid information and outputs the same to the processor core 10.

On the other hand, while the extend instruction is executed, the processor core 10 effects a control in such a manner that the destination valid information outputted from the control circuit 204a in the extended arithmetic unit 20 is stored in the memory stage control register 115a at the rising of the following clock.

In some cases of the pipeline control in the processor core 10, when the destination registers in both the execute stage and memory stage are enabled, and the register number coincides with the register number to be read out in the register read stage, the pipeline may be stalled. When an extended instruction is an instruction that prohibits writing into the register file 105 in the processor core 10, the stall of the pipeline, which is unnecessary at all, can be avoided by providing a function for making the extended arithmetic unit 20 output information indicating prohibition of writing back into the destination register to the processor core 10. In addition, because data is not sent uselessly from the extended arithmetic unit 20 to the processor core 10, power consumption or the like can be saved as well.

In summary, according to the present invention, by providing an extended arithmetic unit instruction as an instruction for controlling the extended arithmetic unit, and a function of controlling execution of the extended arithmetic unit instruction as well as an interface function for connecting the extended arithmetic unit, the extended arithmetic unit can be connected efficiently to the exterior of the processor core. Consequently, by connecting various kinds of extended arithmetic units suitable for application systems to the exterior of the processor core and downsizing

the arrangement of the processor core to the minimum, the system performance can be improved by executing efficient processing between the processor core and extended arithmetic unit while saving the manufacturing costs of the processor core.

- 5 It is to be noted that, besides those already mentioned above, many modifications and variations of the above embodiments may be made without departing from the novel and advantageous features of the present invention. Accordingly, all such modifications and variations are intended to be
10 included within the scope of the appended claims.

What is claimed is:

1. A processor, comprising:

a processor core;

5 a data memory accessed by said processor core; and
an extended arithmetic unit, connected to an exterior
of said processor core, for processing a particular
instruction,

said extended arithmetic unit executing an arithmetic
10 operation by using arithmetic operation data retained in a
register file in said processor core, and outputting a result
of an arithmetic operation directly to said processor core,

said processor core saving the result of the arithmetic
operation executed by said extended arithmetic unit and
15 inputted therefrom in said register file in said processor
core.

2. A processor, comprising:

a processor core;

20 a data memory accessed by said processor core; and
an extended arithmetic unit, connected to an exterior
of said processor core, for processing a particular
instruction,

said processor core, at least including:

25 an instruction memory for storing an instruction to be
executed;

an instruction decode unit for reading out an instruction
from said instruction memory to decode the instruction, in
case that the instruction decoded is an extended arithmetic
30 unit control instruction that should be executed by said
extended arithmetic unit connected to the exterior of said
processor core, said instruction decode unit also outputting
at least an instruction code of said extended arithmetic unit
control instruction to said extended arithmetic unit;

35 a register file for retaining arithmetic operation data
of an arithmetic operation that should be executed by the

instruction decoded, in case that said arithmetic operation data is data of said extended arithmetic unit control instruction, said register file also outputting said arithmetic operation data to said extended arithmetic unit;

5 a first operational section for executing the instruction decoded; and

an extended arithmetic unit, at least including,
a second operational section for executing an arithmetic operation specified by said extended arithmetic unit control instruction by using said arithmetic operation data retained
10 in said register, and outputting an execution result of the arithmetic operation to said processor core.

3. The processor according to claim 1, wherein,
15 in case that the instruction decoded is said extended arithmetic unit control instruction, said processor core outputs to said extended arithmetic unit at least an instruction code that specifies an action involved in an arithmetic operation in said extended arithmetic unit and an
20 instruction valid signal that indicates said instruction code is valid.

4. The processor according to claim 1, wherein said arithmetic operation data outputted to said extended arithmetic unit is
25 a value read out from said register file in said processor core in accordance with a register number specified by a part of said extended arithmetic unit control instruction.

5. The processor according to claim 1, wherein said processor
30 core includes a pipeline control unit for controlling pipeline processing in an interior of said processor core and in said extended arithmetic unit.

6. The processor according to claim 5, wherein said pipeline
35 control unit outputs to said extended arithmetic unit a first pipeline stop signal for suspending execution of an instruction

in said extended arithmetic unit.

7. The processor according to claim 5, wherein said pipeline control unit outputs to said extended arithmetic unit a
5 pipeline flush signal for abandoning execution of an instruction outputted to said extended arithmetic unit.

8. The processor according to claim 5, wherein said pipeline control unit stops execution of an instruction in said
10 processor core in accordance with a second pipeline stop signal for suspending execution of an instruction inputted from said extended arithmetic unit and executed by said processor core.

9. The processor according to claim 1, wherein said extended
15 arithmetic unit outputs to said processor core an arithmetic operation result invalidating signal that invalidates an execution result of an arithmetic operation executed in said processor core.

20 10. The processor according to claim 1, wherein said data memory:

receives from said extended arithmetic unit at least one of an address for memory access, data, a write control signal for controlling data writing, and a read control signal
25 for controlling data reading;

reads out the data from a region specified by said address and outputs the data to said extended arithmetic unit in case that data reading is carried out because said read control signal is asserted; and

30 writes the data inputted from said extended arithmetic unit into a region specified by said address in case that data writing is carried out because said write control signal is asserted.

35 11. The processor according to claim 1, wherein said extended arithmetic unit includes:

a plurality of arithmetic circuits;
a first pipeline register for storing a processing result
by an arithmetic circuit in a preceding stage at a rising of
a following clock; and

5 a second pipeline register for storing a processing
result by an arithmetic circuit in a succeeding stage at the
rising of the following clock.

12. A processor core connected to an extended arithmetic unit
10 for processing a particular instruction to an exterior thereof,
comprising:

an instruction memory for storing an instruction to be
executed;

an instruction decode unit for reading out an instruction
15 from said instruction memory to decode the instruction, in
case that the instruction decoded is an extended arithmetic
unit control instruction that should be executed by said
extended arithmetic unit connected to the exterior of said
processor core, said instruction decode unit also outputting
20 at least an instruction code of said extended arithmetic unit
control instruction to said extended arithmetic unit; and

a register file for retaining arithmetic operation data
of an arithmetic operation that should be executed by the
instruction decoded, and in case that said arithmetic operation
25 data is data for said extended arithmetic unit control
instruction, said register file also outputting said
arithmetic operation data to said extended arithmetic unit
and storing a result of an arithmetic operation executed in
said extended arithmetic unit.

30
13. The processor core according to claim 12, wherein, in
case that the instruction decoded is said extended arithmetic
unit control instruction, said instruction decode unit outputs
to said extended arithmetic unit at least an instruction code
35 that specifies an action involved in an arithmetic operation
by said extended arithmetic unit and an instruction valid

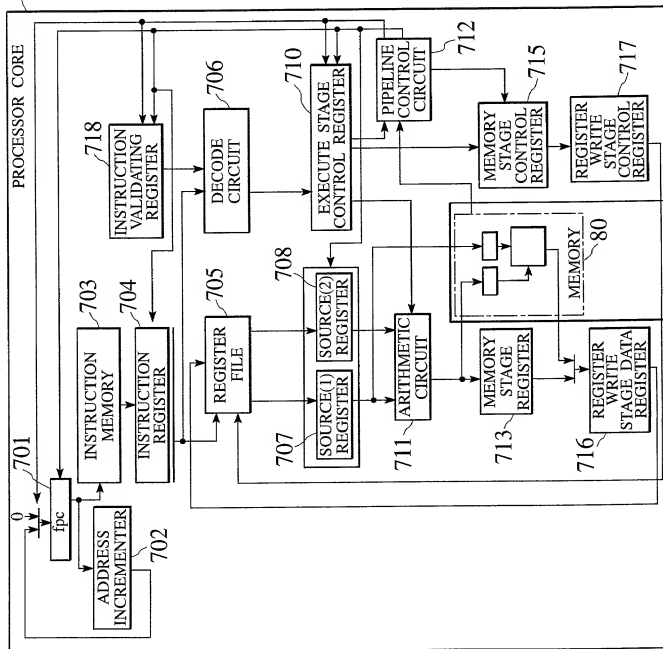
signal that indicates said instruction code is valid.

14. The processor core according to claim 12, wherein said arithmetic operation data outputted to said extended
5 arithmetic unit is a value read out from said register file in said processor core in accordance with a register number specified by a part of said extended arithmetic unit control instruction.
- 10 15. The processor core according to claim 12, further comprising a pipeline control unit for controlling pipeline processing in an internal of said processor core and in said extended arithmetic unit.
- 15 16. The processor core according to claim 15, wherein said pipeline control unit outputs to said extended arithmetic unit a first pipeline stop signal for suspending execution of an instruction in said extended arithmetic unit.
- 20 17. The processor core according to claim 15, wherein said pipeline control unit outputs to said extended arithmetic unit a pipeline flush signal for abandoning execution of an instruction outputted to said extended arithmetic unit.
- 25 18. The processor core according to claim 15, wherein said pipeline control unit stops execution of an instruction in said processor core in accordance with a second pipeline stop signal, inputted from said extended arithmetic unit, for suspending execution of an instruction in said processor core.
- 30 19. The processor core according to claim 15, wherein said pipeline control unit receives from said extended arithmetic unit an arithmetic operation result invalidating signal that invalidates an execution result of an arithmetic operation
35 in said processor core, and invalidates the execution result of the arithmetic operation in said processor core.

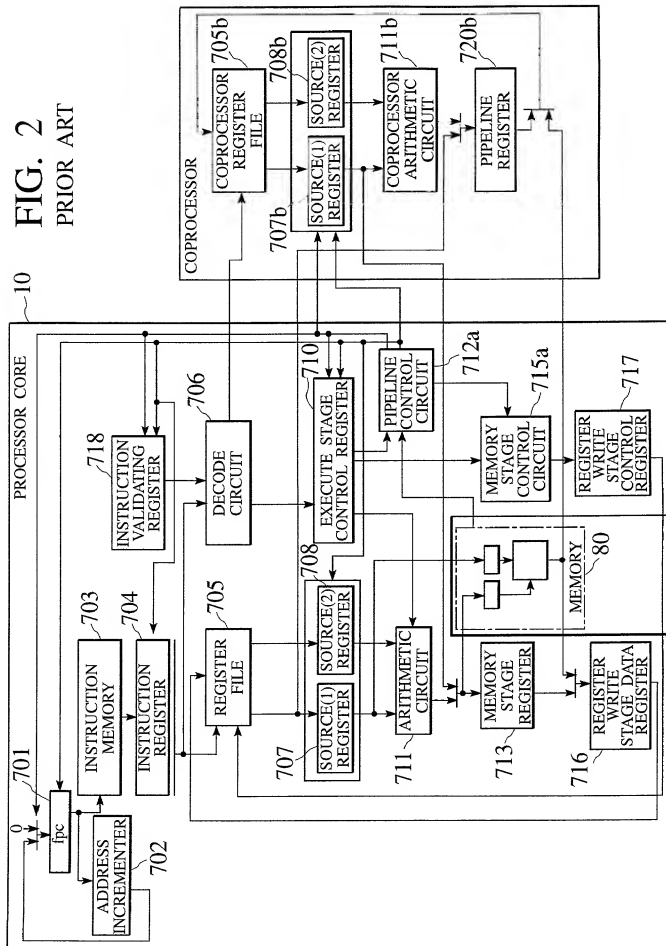
ABSTRACT OF THE DISCLOSURE

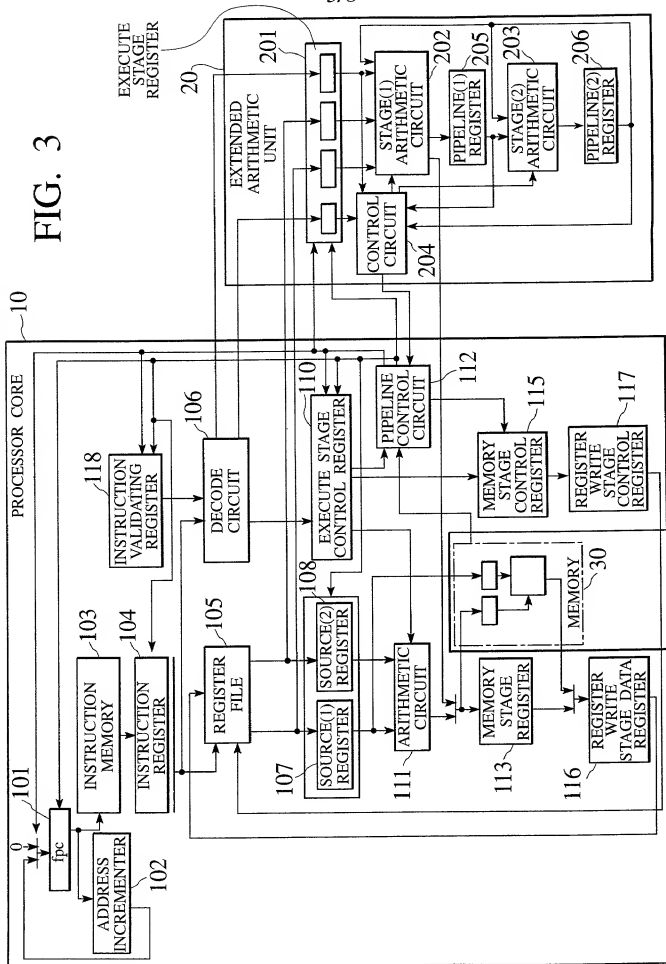
A processor core for realizing efficient operation processing by connecting an extended arithmetic unit to its exterior and a processor incorporating such a processing core are provided. The processor includes the processor core, a data memory accessed by the processor core, and the extended arithmetic unit connected to the exterior of the processor core for processing a particular instruction. The extended arithmetic unit executes an arithmetic operation by using arithmetic operation data retained in a register file in the processor core, and directly outputs an arithmetic operation result to the processor core. Then, the processor core saves the result of the arithmetic operation executed by the extended arithmetic unit and inputted therefrom in the register file in the processor core.

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FIG. 1
PRIOR ART

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FIG. 2
PRIOR ART



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FIG. 4

	31	28 27	24 23	20 19	16 15	0
add dst, src1, src2	dst	src1	src2	0000	0000	0000 0000
	dst	src1	0000	0001	0000	0000 0000
load dst, (src1)	0000	src1	src2	0010	0000	0000 0000
store (src1), src2	0000	src1	0000	0011	0000	0000 0000
jump src1	dst	src1	src2	0100	code16	
ext dst, src1, src2, code16						

031260" 95499960

FIG. 6

ALU CONTROL INPUT	ALU OUTPUT
00 (nop)	SOURCE(1) REGISTER
01 (add)	SOURCE(1) REGISTER + SOURCE(2) REGISTER
10 (ext)	INPUT FROM EXTENDED ARITHMETIC UNIT

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FIG. 7

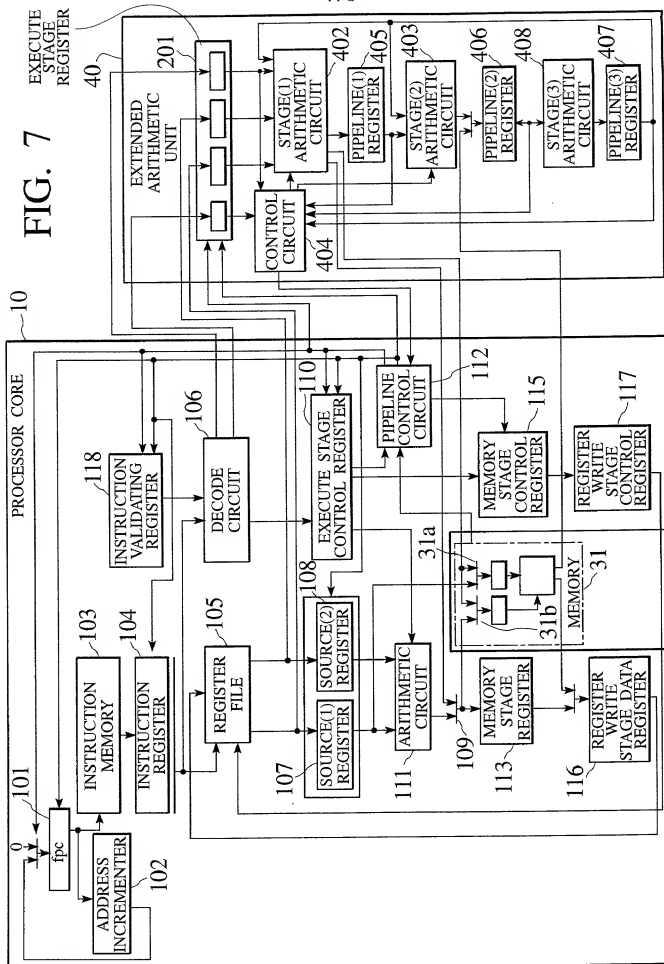
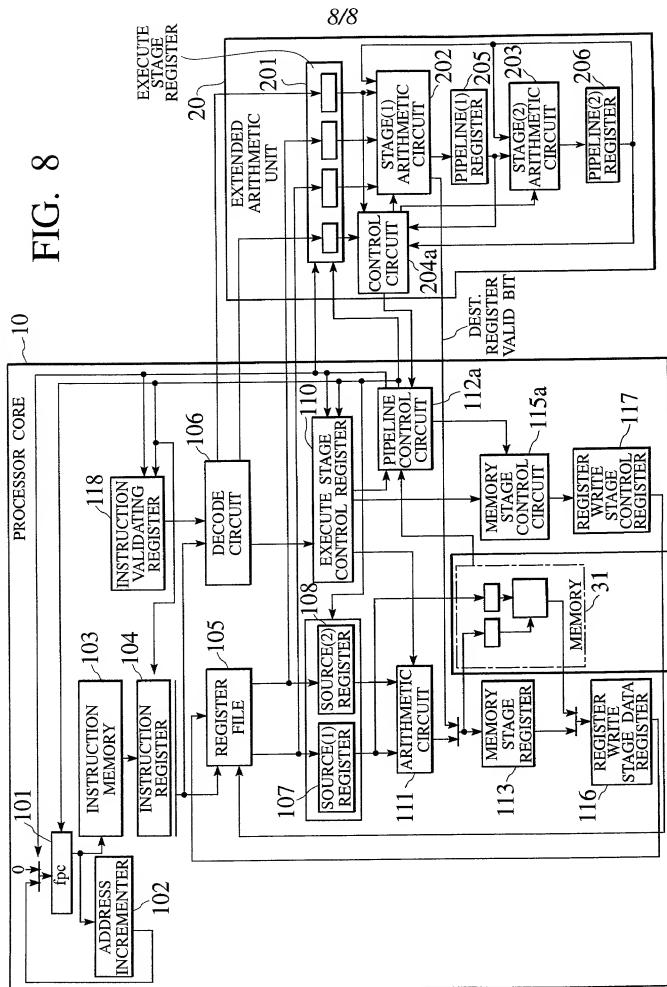


FIG. 8



466 2.1.7.3

We (I) hereby claim the benefit under Title 35, United States Code, § 119(e) of any United States provisional application(s) listed below.

(Application Number)

(Filing Date)

(Application Number)

(Filing Date)

We (I) hereby claim the benefit under 35 U.S.C. § 120 of any United States application(s), or § 365(c) of any PCT International application designating the United States, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States or PCT International application in the manner provided by the first paragraph of 35 U.S.C. § 112, I acknowledge the duty to disclose information which is material to patentability as defined in 37 CFR § 1.56 which became available between the filing date of the prior application and the national or PCT International filing date of this application.

Application Serial No.

Filing Date

Status (pending, patented,
abandoned)

And we (I) hereby appoint: Norman F. Oblon, Registration Number 24,618; Marvin J. Spivak, Registration Number 24,913; C. Irvin McClelland, Registration Number 21,124; Gregory J. Maier, Registration Number 25,599; Arthur I. Neustadt, Registration Number 24,854; Richard D. Kelly, Registration Number 27,757; James D. Hamilton, Registration Number 28,421; Eckhard H. Kuesters, Registration Number 28,870; Robert T. Pous, Registration Number 29,099; Charles L. Gholz, Registration Number 26,395; Vincent J. Sunderdick, Registration Number 29,004; William E. Beaumont, Registration Number 30,996; Steven B. Kelber, Registration Number 30,073; Robert F. Gnuse, Registration Number 27,295; Jean-Paul Lavalleye, Registration Number 31,451; Timothy R. Schwartz, Registration Number 32,171; Stephen G. Baxter, Registration Number 32,884; Martin M. Zoltick, Registration Number 35,745; Robert W. Hahl, Registration Number 33,893; Richard L. Treanor, Registration Number 36,379; Steven P. Weihrouch, Registration Number 32,829; John T. Goolkasian, Registration Number 26,142; Marc R. Labgold, Registration Number 34,651; William J. Healey, Registration Number 36,160; Richard L. Chinn, Registration Number 34,305; Steven E. Lipman, Registration Number 30,011; Carl E. Schlier, Registration Number 34,426; James J. Kulbaski, Registration Number 34,648; Catherine B. Richardson, Registration Number 39,007; Richard A. Neifeld, Registration Number 35,299; J. Derek Mason, Registration Number 35,270; and Jacques M. Dulin, Registration Number 24,067; our (my) attorneys, with full powers of substitution and revocation, to prosecute this application and to transact all business in the Patent Office connected therewith; and we (I) hereby request that all correspondence regarding this application be sent to the firm of OBLON, SPIVAK, MCCLELLAND, MAIER & NEUSTADT, P.C., whose Post Office Address is: Fourth Floor, 1755 Jefferson Davis Highway, Arlington, Virginia 22202.

We (I) declare that all statements made herein of our (my) own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issuing thereon.

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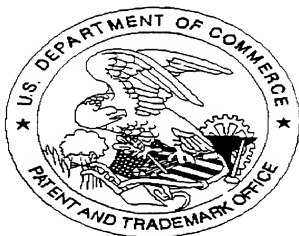
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